# Experimental study of approximation algorithms for a graph clustering problem

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# Clustering problems

In the *clustering problem* we must split a set of objects into several subsets based on the similarity of the objects to each other. A set of objects can be represented as vertices of a graph, and the similarity of objects can be specified as edges of this graph.

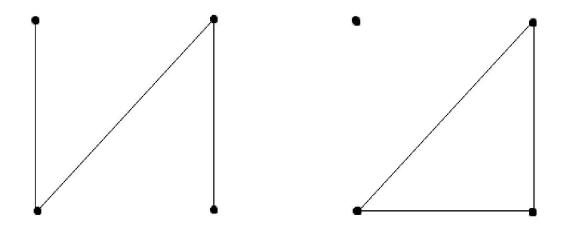
A graph is clustered if each of its components is a clique (*cluster*). Clustering problems are related to *unsupervised learning*. However, *semi-supervised* methods and algorithms are also applicable. In this approach, we have a supervisor who can partition some objects across clusters.

# **Basic definitions**

The distance  $\rho(G_1, G_2)$  between two labeled graphs  $G_1 = (V, E_1)$  and  $G_2 = (V, E_2)$  is the cardinality of the symmetric difference  $E_1 \Delta E_2$  (the elements of this set are called *disagreements*).

For the vertex *v* of the graph G = (V, E) we call a *neighborhood*  $N_G(v)$  the set of  $u \in V$  joined with *v*.

# Cluster graph example



We have to remove one edge and add one edge.

# Semi supervised graph clustering

**k-SEMI-SUPERVISED GRAPH CLUSTERING.** The input is a graph G = (V, E), an integer  $2 \le k \le |V|$ , a set  $Z = \{z_1, ..., z_k\} \subset V$ . The aim is to find cluster graph *C* with *k* clusters minimizing the number of disagreements. Additionally, vertices of *Z* must belong to different clusters of *C*.

The problem is *NP*-hard.

# 3-approximation algorithm for 2-ss-clustering

#### The Neighborhood Semi-Supervised Algorithm (NS).

Construct the set *F* of feasible solutions according to the rules.

- a) For each vertex  $v \in V \setminus \{z_1, z_2\}$  build two cluster graphs  $C_v^{-1}$  and  $C_v^{-2}$  by the following. The 1st cluster of the 1st graph is  $\{v\} \cup (N_G(v) \setminus \{z_1\}) \cup \{z_2\}$ , the 1st cluster of the 2nd graph is  $\{v\} \cup (N_G(v) \setminus \{z_2\}) \cup \{z_1\}$ . Both 2nd clusters contain not yet clustered vertices.
- b) For each vertex  $v \in \{z_1, z_2\}$  build cluster graphs  $C_v$  by the following. The 1st cluster is  $(\{v\} \cup N_G(v)) \setminus \{x\}$ . 2nd cluster contains not yet clustered vertices ( $x = z_1$  if  $v = z_2$ ,  $x = z_2$  if  $v = z_1$ ).

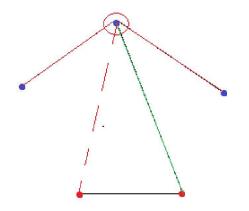
Return  $C_{NS} \in F$  with minimum  $\rho(G, C_{NS})$ .

## Local search for 2-ss-clustering

#### LocalSearch.

Let *C* be a cluster graph for a graph G = (V, E). For each  $v \in V \setminus \{z_1, z_2\}$  let  $C_v$  be the same cluster graph as *C*, except with *v* in the opposite cluster. We then define  $\lambda_v = \rho(G, C) - \rho(G, C_v)$ , the improvement caused by the change. Let *u* be the vertex with maximum  $\lambda_v$ . If  $\lambda_v \leq 0$ , stop, otherwise let  $C \leftarrow C_w$  and repeat.

### Local search for 2-ss-clustering



When we move a blue vertex to a red cluster, we don't count the green edge, but we count the red ones. Therefore, the value of the objective function will increase by 2.

# Approximation algorithms for 2-ss-clustering

The 2-approximation Neighborhood Semi-Supervised with LocalSearch Algorithm (NSLS) from is the application of local search to each cluster graph  $C \in F$ . Let  $C_{NSLS}$  be a solution built by this algorithm.

We also want to research **Pre-Clustered Neighborhood Semi-Supervised with LocalSearch Algorithm (PNSLS)** which applies local search only for  $C_{z1}$  and  $C_{z2}$  from *F*. Let  $C_{PNSLS}$  be a solution built by this algorithm.

Complexity of NS -  $O(n^2)$ , NSLS -  $O(n^4)$ , PNSLS -  $O(n^3)$ .

## Average working time of NS, NSLS and PNSLS

n	500	600	700	800	900	1000	1200	1400	1600	1800	2000	2500	3000
NS	0	0	0	0	0	1	1	2.21	4.01	5.99	8.04	15.49	27.65
NSLS	0.07	1	1.04	2	3	4	6.65	10.73	16.45	25.32	32.84	66.45	119.39
PNSLS	0	0	0	0	0	1	1	2.32	4.02	6	8.04	15.78	27.78

# Background of experimental study

It's easy to see that  $\rho(G, C_{NS}) \leq \rho(G, C_{NSLS})$  and  $\rho(G, C_{PNSLS}) \leq \rho(G, C_{NSLS})$  for any G = (V, E). Let's define  $E_{NS}(G) = \rho(G, C_{NS}) / \rho(G, C_{NSLS})$  and  $E_{PNSLS}(G) = \rho(G, C_{PNSLS}) / \rho(G, C_{NSLS})$  as errors of **NS** and **PNSLS** relatively to **NSLS**.

Then define  $E_{NS}(n)$  and  $E_{PNSLS}(n)$  as expected values of the  $E_{NS}$  and  $E_{PNSLS}$  for all graphs with *n* vertices. Let's formulate the main assumption to be investigated.

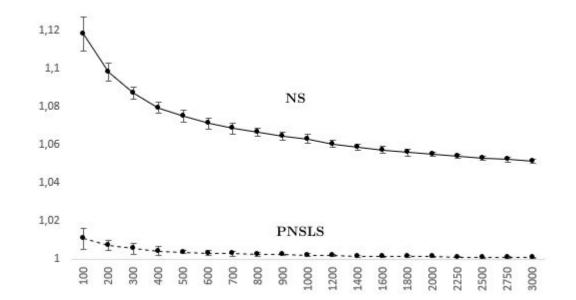
**Assumption.** As the number of vertices *n* increases,  $E_{PNSLS}(n)$  tends to 1 and  $E_{NS}(n)$  doesn't tend to 1.

# **Experimental study**

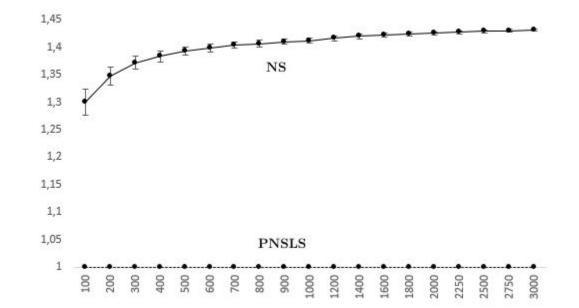
The experimental study was done on random graphs G(n, p) generated with *Erdős-Rényi model*. The number of vertices  $n \in \{100, 200, ..., 3000\}$ , graph density  $p \in \{0.33, 0.5, 0.67\}$ . 100 examples were solved for each *n* and *p*. Based on the sample data, the sample mean of  $E_{NS}(n)$  and  $E_{PNSLS}(n)$  was calculated. Further, with significance level  $\alpha = 0.05$ , we calculated the confidence interval. We used the quantile of the normal distribution to calculate the confidence interval.

Statistical validity was obtained by the *Kolmogorov-Smirnov test*. For each *n* and *p*, the statistic was less than the critical value 1.36.

# Error plot of $\boldsymbol{E}_{NS}$ and $\boldsymbol{E}_{PNSLS}$ for p = 0.33



# Error plot of $\boldsymbol{E}_{NS}$ and $\boldsymbol{E}_{PNSLS}$ for p = 0.67



# Thank you for your attention.